# SHORT PROOF THAT KNESER GRAPHS ARE HAMILTONIAN FOR $n \geqslant 4 k$ 

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#### Abstract

For integers $n \geqslant k \geqslant 1$, the Kneser graph $K(n, k)$ is the graph with vertex set $V=[n]^{(k)}$ and edge set $E=\left\{\{x, y\} \in V^{(2)}: x \cap y=\varnothing\right\}$. Chen proved that for $n \geqslant 3 k$, Kneser graphs are Hamiltonian and later improved this to $n \geqslant 2.62 k+1$. Furthermore, Chen and Füredi gave a short proof that if $k \mid n$, Kneser graphs are Hamiltonian for $n \geqslant 3 k$. In this note, we present a short proof that does not need the divisibility condition, i.e., we give a short proof that $K(n, k)$ is Hamiltonian for $n \geqslant 4 k$.


## §1. Introduction

Throughout the paper, let $n \geqslant k \geqslant 1$ be integers and set $[n]=\{1, \ldots, n\}$. For a set $A$, define $A^{(k)}$ to be the set of all $k$-element subsets (or $k$-subsets) of $A$. The Kneser graph $K(n, k)$ has vertex set $[n]^{(k)}$ and two vertices form an edge if and only if they are disjoint (as subsets of [ $n$ ]). With rather involved proofs Chen [2] showed that Kneser graphs with $n$ linear in $k$, and even triangle-free Kneser graphs [3] contain Hamiltonian cycles. More precisely, in [3] she showed the following.

Theorem 1.1. If $n \geqslant 2.62 k+1$, then $K(n, k)$ is Hamiltonian.
Chen and Füredi [4] simplified the proof for the case when $k \mid n$ and $n \geqslant 3 k$. Katona [5] conjectured that, apart from finitely many exceptions, $K(n, k)$ is Hamiltonian if $n \geqslant 2 k+1$. Recently Mütze, Nummenpalo, and Walczak [6] showed that for $k \geqslant 3$, the Kneser graph $K(2 k+1, k)$ is Hamiltonian (and they also provide a more exhaustive coverage of the previous work in this area).

In this note, we elaborate the short proof due to Chen and Füredi to work for the general case by removing the divisibility condition. More precisely, we give a short proof that

Theorem 1.2. $K(n, k)$ is Hamiltonian for $n \geqslant 4 k$.
A Gray-Code is an enumeration $x_{1} \ldots x_{m}$ of all sets in $[n]^{(k)}$ such that each two consecutive sets and in addition $x_{m}, x_{1}$ differ by exactly one element (of $[n]$; in general, we say two $k$ sets $x, y \in[n]^{(k)}$ differ by $i$ elements if $|x \backslash y|=i$. The existence of Gray-Codes follows

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easily by induction (see e.g., [7]) and they were also used in [4]. Observing that the edges of $K_{n}^{(k)}$ correspond to the vertices of $K(n, k)$ and a matching of size $s$ in $K_{n}^{(k)}$ corresponds to a clique of size $s$ in $K(n, k)$, we get the following corollary to Baranyai's theorem [1].

Theorem 1.3. Let $n \geqslant k$ and $a_{1}, \ldots, a_{t} \leqslant \frac{n}{k}$ be integers such that $\sum_{i=1}^{t} a_{i}=\binom{n}{k}$. Then $K(n, k)$ can be partitioned into cliques $A_{i}$ with $\left|A_{i}\right|=a_{i}$ for $i \in[t]$.

## §2. Short proof of Theorem 1.2

For clarity, we first give the proof for $n \geqslant 5 k$ and afterwards, in Remark 2.1, we will go through the proof again inserting the additional arguments for $n \geqslant 4 k$.

Proof of Theorem 1.2 if $n \geqslant 5 k$. Let $p, r, m, q \in \mathbb{Z}$ such that $n=p k+r$ with $0 \leqslant r \leqslant k-1$ and $\binom{n}{k}=(m-1) p+q$ with $1 \leqslant q \leqslant p$. Set $b=b(q)=\max \{4-q, 0\}$ and define

$$
a_{i}= \begin{cases}p & \text { for } i \in[m-1-b]  \tag{2.1}\\ p-1 & \text { for } i \in\{m-b, \ldots, m-1\} \\ q+b & \text { for } i=m\end{cases}
$$

Then we have $\sum_{i=1}^{m} a_{i}=\binom{n}{k}$ and Theorem 1.3 provides a partition of $K(n, k)$ into cliques $A_{1}, \ldots, A_{m}$ with $\left|A_{i}\right|=a_{i} \geqslant 4$ for all $i \in[m]$. For $i \in[m]$, define marking vertices as follows. If there is an $x_{i} \in A_{i}$ with $n \in x_{i}$, set $x_{i}$ to be the marking vertex of $A_{i}$. If there is no vertex containing $n$ in $A_{i}$, we choose an arbitrary vertex $x_{i} \in A_{i}$ as marking vertex of $A_{i}$. We call the set of marking vertices $M$ and note that $M$ contains $M^{\prime}=\left\{z \cup\{n\}: z \in[n-1]^{(k-1)}\right\}$. Next, we use a Gray-Code on $[n-1]^{(k-1)}$ to obtain an enumeration $x_{1}^{\prime} \ldots x_{m^{\prime}}^{\prime}$ of $M^{\prime} \subseteq M$ with $\left|x_{i}^{\prime} \backslash x_{i+1}^{\prime}\right|=1$ for all $i \in \mathbb{Z} / m^{\prime} \mathbb{Z}$. Further, we consider a map $\varphi: M \backslash M^{\prime} \rightarrow M^{\prime}$, so that for each $x \in M \backslash M^{\prime}$, we have $|\varphi(x) \backslash x|=1$ (this is possible since for all $a \in x \in M \backslash M^{\prime}$, the vertex $x \backslash\{a\} \cup\{n\}$ is in $\left.M^{\prime}\right)$. Thus, the enumeration

$$
x_{1}^{\prime} \varphi^{-1}\left(x_{1}^{\prime}\right) x_{2}^{\prime} \varphi^{-1}\left(x_{2}^{\prime}\right) \ldots x_{m^{\prime}}^{\prime} \varphi^{-1}\left(x_{m^{\prime}}^{\prime}\right)=y_{1} \ldots y_{m}
$$

of $M$ (here $\varphi^{-1}\left(x_{i}^{\prime}\right)$ stands for an arbitrary enumeration of $\left.\varphi^{-1}\left(x_{i}^{\prime}\right)\right)$ has the property that $y_{i}$ and $y_{i+1}$ differ by at most two elements (for $i \in \mathbb{Z} / m \mathbb{Z}$ ). Since $\left|A_{i}\right| \geqslant 4$, this yields that there is a vertex $z_{i} \in A\left(y_{i}\right)$ which is disjoint to $y_{i+1}$, where $A\left(y_{i}\right)$ is the clique among $A_{1}, \ldots, A_{m}$ that contains $y_{i}$. Thus, denoting by $\alpha_{i}$ an enumeration of $A\left(y_{i}\right)$ that starts with $y_{i}$ and ends with $z_{i}$, we get that $\alpha_{1} \ldots \alpha_{m}$ is a Hamiltonian cycle.

Remark 2.1. Here we mention the modifications that let the proof above work for all $n \geqslant 4 k$. Note that for $k=1$ the result is trivial, so assume $k \geqslant 2$. First, using 0 instead of $b$ in (2.1) yields that $\left|A_{i}\right| \geqslant 4$ for $i \in[m-1]$ and $\left|A_{m}\right|=q$. If $q \geqslant 4$, the same proof
as above still works. So we can assume that $q \leqslant 3$ and hence, there is an element of $[n]$, w.l.o.g. $n$, that is not contained in any vertex of $A_{m}$. For $i \in[m-1]$, we define marking vertices as before and for $A_{m}$, we do not define a marking vertex. Proceeding as above gives an enumeration $y_{1} \ldots y_{m-1}$ of the marking vertices with the property that $y_{i}$ and $y_{i+1}$ differ by at most two elements (for $i \in \mathbb{Z} /(m-1) \mathbb{Z}$ ) and so we still know that the vertices $z_{i}$ exist as before (for $i \in[m-1]$ ).

Thus, we get as above that $\alpha_{1} \ldots \alpha_{m-1}$ is a cycle $C$ which covers all but at most three vertices $v_{1}, v_{2}, v_{3} \in A_{m}$. Note that for each $v_{i}$, we can choose a marking vertex $y_{j(i)} \in M^{\prime}$ with $\left|v_{i} \backslash y_{j(i)}\right|=1$. Since $k \geqslant 2$ and $A_{m}$ is a clique, $y_{j(i)} \neq y_{j\left(i^{\prime}\right)}$ whenever $v_{i} \neq v_{i^{\prime}}$. Further, $\left|A\left(y_{j(i)}\right)\right| \geqslant 4$ implies that there are two vertices $u_{i}^{1}, u_{i}^{2} \in A\left(y_{j(i)}\right)$ which are disjoint to $v_{i}$. Note, that the enumeration $\alpha_{j(i)}$ of $A\left(y_{j(i)}\right)$ was arbitrary apart from the start $\left(y_{j(i)}\right)$ and the end $\left(z_{i}\right)$. So we can additionally request that $u_{i}^{1}$ and $u_{i}^{2}$ are next to each other in this enumeration and insert $v_{i}$ into $C$ between $u_{i}^{1}$ and $u_{i}^{2}$, obtaining a Hamiltonian cycle.

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