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Topological circles and Euler tours in locally finite graphs

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#### Abstract

We obtain three results concerning the topological cycle space  $\mathcal{C}$  of a locally finite connected graph G. Confirming a conjecture of Diestel we show that through every edge set  $E \in \mathcal{C}$  there is a topological Euler tour, a continuous map from the circle  $S^1$  to the end compactification |G| of G that traverses every edge in E exactly once and traverses no other edge.

Second, we show that for every sequence  $(\tau_i)_{i \in \mathbb{N}}$  of topological x-y paths in |G| there is a topological x-y path in |G| all of whose edges lie eventually in every member of some fixed subsequence of  $(\tau_i)$ . As a corollary we obtain a short proof of one of the fundamental theorems about C, that every  $E \in C$  is a disjoint union of edge sets of circles in |G|.

Third, we show that every set of edges not containing a finite odd cut of G extends to an element of C.

## 1 Introduction

Erdős et al. [16] characterised the infinite graphs G containing an *Euler tour*, that is, a two-way infinite walk traversing every edge of G precisely once. Although their result is best possible, it is not really satisfying: graphs with more than two ends cannot have such an Euler tour, so we cannot generalise to arbitrary infinite graphs the well-known theorems about Euler tours in finite ones. However, Diestel and Kühn [13] proposed a new concept of an Euler tour, called a *topological Euler tour*, that does allow such theorems to generalise to arbitrary infinite graphs, at least locally finite ones.

A topological Euler tour of a locally finite (multi-)graph G is a continuous map  $\sigma : S^1 \to |G|$  traversing every edge of G precisely once. Here, |G| is the Freudenthal compactification [17] of G, a topological space consisting of G and its ends, and  $S^1$  is the unit circle in the real plane. (See Section 2 for precise definitions.) Unlike the approach of [16], this definition makes it possible to have a topological Euler tour in a graph that has more than two ends.

For example, the following theorem can be considered as a generalisation of Euler's theorem that a finite connected multigraph has an Euler tour if and only if every vertex has even degree.

**Theorem 1.1** ([13]). Let G be a connected locally finite multigraph. Then, G admits a topological Euler tour if and only if every finite cut of G is even.

The concept of a topological Euler tour and Theorem 1.1 are part of a far-reaching research project initiated by Diestel that uses topological concepts to generalise some classical parts of finite graph theory to infinite graphs; see [10] or [11, Chapter 8.5] for an introduction. In this context, a *circle* in |G| is a homeomorphic image of  $S^1$ . The *(topological) cycle space* C(G) of G is a vector space over  $\mathbb{Z}_2$  built from edge sets of circles in |G|; see Section 2. This concept has allowed the generalisation to infinite graphs of most of the well-known theorems about the cycle space of a finite graph, most of which fail for infinite graphs if the usual finitary notion of the cycle space is applied [2, 3, 4, 5, 13, 20].

A result in this context, which is used in the proof of Theorem 1.1, is the following.

**Theorem 1.2** ([13]). Let G be a connected locally finite multigraph. Then,  $E(G) \in C(G)$  if and only if G admits a topological Euler tour.

If  $H \subseteq G$  are finite connected multigraphs and  $E(H) \in \mathcal{C}(G)$  then it follows easily from (the finite version of) Theorem 1.2 that H admits an Euler tour. Subgraphs of infinite multigraphs however are harder to handle, and one reason why this is the case is that, in general, they have a different end-topology compared to the original graph. It is thus not surprising that although Theorem 1.2 was one of the first results in this research project, it still remained open to prove

**Theorem 1.3.** Let  $H \subseteq G$  be locally finite multigraphs such that the closure  $\overline{H}$  of H in |G| is topologically connected. Then  $E(H) \in C(G)$  if and only if H admits a topological Euler tour in |G|.

Here a topological Euler tour of H in |G| is a continuous map from  $S^1$  to |G| traversing each edge of H precisely once and meeting no other edge at an inner point. Theorem 1.3 was conjectured by Diestel [9], and it is the main result of this paper. We prove it in Section 4.

Our next result is a lemma whose main idea has already appeared implicitely several times (e.g. in [20] and in [18, Theorem 4]) and which could have further applications. Given a locally finite graph G and two vertices x, y in G, this lemma facilitates the construction of a topological x-y path contained in a fixed subspace of |G|. Unlike finite paths, such a topological path cannot be constructed step by step attaching one edge at a time, because the structure of a topological path can be arbitrarily complicated: in general, a topological path consists of double rays whose ordering in the path can be of any countable ordinal type (see for example the circle in Figure 1, in which the double rays are arranged like the rational numbers). The main idea behind our lemma is to construct such a topological path as a limit of x-y paths in (finite) subgraphs of G. It states that any infinite sequence of topological paths between two fixed vertices x, y in |G| has a subsequence whose limit contains a topological x-y path. More precisely, given a sequence  $E_1, E_2, \ldots$  of sets, let us write

$$\liminf(E_n) := \bigcup_{i \in \mathbb{N}} \bigcap_{j > i} E_j$$

for the set of elements eventually in  $E_n$ . For a subspace X of |G| we write E(X) for the set of edges contained in X. Similarly, for a topological path  $\tau$  we write  $E(\tau)$  for the set of edges contained in the image of  $\tau$ .

**Lemma 1.4.** Let G be a locally finite graph, let  $x, y \in V(G)$  and let  $(\tau_i)_{i \in \mathbb{N}}$ be a sequence of topological x-y paths in |G|. Then, there is an infinite subsequence  $(\tau_{a_i})_{i \in \mathbb{N}}$  of  $(\tau_i)$  and a topological x-y path  $\sigma$  in |G| such that  $E(\sigma) \subseteq$  $\liminf(E(\tau_{a_i}))$ . Moreover, if no  $\tau_i$  traverses any edge more than once then  $E(G) \setminus E(\sigma) \subseteq \liminf(E(G) \setminus E(\tau_{a_i}))$ , no edge is traversed by  $\sigma$  more than once, and for every finite subset F of  $E(\sigma)$  there is an  $n \in \mathbb{N}$  such that the linear ordering of F induced by  $\sigma$  coincides with that induced by  $\tau_{a_i}$  for every  $a_i > n$ .

In Section 5 we will use Lemma 1.4 in order to obtain an elementary proof of the following result, which is one of the most important tools in this research project:

**Theorem 1.5** ([14]). For any locally finite graph G, every element of C(G) is a disjoint union of edge-sets of circles.

One of the applications of Theorem 1.5 is in the proof of Theorem 1.2. Theorem 1.5 has a long history; its first proof, in [14], applied to non-locally-finite graphs as well and was quite technical. A simpler proof for locally finite graphs inspired by [25] appears in [11, Theorem 8.5.8], but this proof still uses a non-trivial fact from topology and another result of Diestel and Kühn [15], that if G is locally finite and  $X \subseteq |G|$  is closed and connected then X is path-connected. Our proof is elementary and based only on Lemma 1.4.

Having seen Theorem 1.1 and Theorem 1.3 it is natural to ask for a generalisation to infinite multigraphs of the following result of Jaeger. For a multigraph G and  $H \subseteq E(G)$ , we say that H extends to an element of  $\mathcal{C}(G)$  if there is  $H' \in \mathcal{C}(G)$  with  $H \subseteq H'$ .

**Theorem 1.6** ([22]). If G is a finite multigraph and  $H \subseteq E(G)$  then H extends to an element of  $\mathcal{C}(G)$  if and only if H contains no cut of G of odd cardinality.

In Section 6 we prove this assertion for locally finite G and discuss a possible application to hamiltonicity problems.

# 2 Definitions

Unless otherwise stated, we will be using the terminology of [11] for graph theoretical concepts and that of [1] for topological ones. Let G = (V, E) be a *locally finite* multigraph — i.e. every vertex has a finite degree — fixed throughout this section.

A 1-way infinite path is called a ray, a 2-way infinite path is a double ray. A tail of the ray R is a final subpath of R. Two rays R, L in G are equivalent if no finite set of vertices separates them. The corresponding equivalence classes of rays are the ends of G. We denote the set of ends of G by  $\Omega = \Omega(G)$ .

Let G bear the topology of a 1-complex<sup>1</sup>. To extend this topology to  $\Omega$ , let us define for each end  $\omega \in \Omega$  a basis of open neighbourhoods. Given any finite set  $S \subset V$ , let  $C = C_G(S, \omega)$ , or just  $C(S, \omega)$  if G is fixed, denote the

<sup>&</sup>lt;sup>1</sup>Every edge is homeomorphic to the real interval [0, 1], the basic open sets around an inner point being just the open intervals on the edge. The basic open neighbourhoods of a vertex x are the unions of half-open intervals [x, z), one from every edge [x, y] at x.

component of G - S that contains some (and hence a tail of every) ray in  $\omega$ , and let  $\Omega(S, \omega)$  denote the set of all ends of G with a ray in C. As our basis of open neighbourhoods of  $\omega$  we now take all sets of the form

$$C(S,\omega) \cup \Omega(S,\omega) \cup E'(S,\omega) \tag{1}$$

where S ranges over the finite subsets of V and  $E'(S, \omega)$  is any union of halfedges (z, y], one for every S-C edge e = xy of G, with z an inner point of e. Let |G| denote the topological space of  $G \cup \Omega$  endowed with the topology generated by the open sets of the form (1) together with those of the 1-complex G.

It can be proved (see [12]) that in fact |G| is the Freudenthal compactification [17] of the 1-complex G.

A circle in |G| is the image of a homeomorphism from  $S^1$ , the unit circle in  $\mathbb{R}^2$ , to |G|. An arc in |G| is a homeomorphic image of the real interval [0, 1] in |G|.

A subset D of E is a *circuit* if there is a circle C in |G| such that  $D = \{e \in E | e \subseteq C\}$ . Call a family  $(D_i)_{i \in I}$  of subsets of E thin, if no edge lies in  $D_i$  for infinitely many i. Let the sum  $\sum_{i \in I} D_i$  of this family be the set of all edges that lie in  $D_i$  for an odd number of indices i, and let the cycle space C(G) of G be the set of all sums of (thin families of) circuits.

A topological Euler tour of G is a continuous map  $\sigma : S^1 \to |G|$  such that every inner point of an edge of G is the image of exactly one point of  $S^1$  (thus, every edge is traversed exactly once, and in a "straight" manner).

A *trail* in a multigraph G is a walk in G that traverses no edge more than once.

If H is a finite submultigraph of the locally finite multigraph G, then let  $H^*$ denote the multigraph obtained from G by contracting each component C of G - H into a single vertex  $u_C$ , deleting loops but keeping multiple edges. If Xis a topological path in |G| then let  $X \upharpoonright H$  be the walk induced in  $H^*$  by X(induced by replacing each maximal subpath of X in a component C of G - Hby  $u_C$ ). If  $X \upharpoonright H$  traverses some edge more than once then pick a trail with edges in  $E(X \upharpoonright H)$  between the endvertices of  $X \upharpoonright H$  and denote it by X|H; if  $X \upharpoonright H$  is a trail then let  $X|H = X \upharpoonright H$ .

#### **3** Basic facts

The following basic fact can be found in [21, p. 208].

**Lemma 3.1.** The image of a topological path with endpoints x, y in a Hausdorff space X contains an arc in X between x and y.

The well known orthogonality relation between elements of the cycle space and cuts has also been extended to locally finite graphs and provides a very useful tool [11, Theorem 8.5.8]:

**Lemma 3.2.** Let  $F \subseteq E(G)$ . Then  $F \in C(G)$  if and only if F meets every finite cut in an even number of edges.

The following lemma, proved e.g. in [11, Theorem 8.5.5], says that an arc is not allowed to jump a finite cut:

**Lemma 3.3.** If F = E(X, Y) is a finite cut of G (where  $\{X, Y\}$  is a bipartition of V(G)) then any arc in |G| with one endpoint in X and one in Y contains an edge from F.

### 4 Euler tours

In this section we prove Theorem 1.3, but before we do so let us consider an example indicating why, contrary to the case when G is finite, this is harder to prove than Theorem 1.2. Let me first sketch the proof of Theorem 1.2 [13]. The harder implication is to show that if  $E(G) \in \mathcal{C}(G)$  then G has a topological Euler tour. To prove this, Diestel & Kühn [13] first applied Theorem 1.5 to obtain a decomposition of E(G) into a set  $\mathcal{B}$  of pairwise edge-disjoint circuits. Then, they picked one of those circuits  $B_0$  and a homeomorphism  $\sigma_0$  from  $S^1$  to the defining circle of  $B_0$ . They proceeded inductively in infinitely many steps, in each step n considering an elements B of  $\mathcal{B}$  such that B is incident with a vertex v visited by  $\sigma_{n-1}$  but the edges of B are not traversed by  $\sigma_{n-1}$ . Then, they modified  $\sigma_{n-1}$  locally in order to obtain a new continuous map  $\sigma_n$  that traverses B as well as all elements of  $\mathcal{B}$  traversed by  $\sigma_{n-1}$ . This process yielded a limit map  $\sigma$  whose image contained every edge exactly once, and the remaining task was to show that this map was continuous.

Now consider the graph F of Figure 1, a well known object from [13, 10]. Formally, it is defined as follows. Let V be the set of finite 01 sequences, including the empty sequence  $\emptyset$ . Define a tree on V by joining every  $\ell \in V$  to its two one-digit extensions, the sequences  $\ell 0$  and  $\ell 1$ . For every  $\ell \in V$ , add another edge  $e_{\ell}$  between the vertices  $\ell 01$  and  $\ell 10$ , and let  $D_{\ell}$  denote the double ray consisting of  $e_{\ell}$  and the two rays starting at  $e_{\ell}$  whose vertices have the form  $\ell 1000 \dots$  and  $\ell 0111 \dots$  Finally, let L be the double ray whose vertices are  $\emptyset$ and the all-zero and the all-one sequences, and let  $\mathcal{D} := \{L\} \cup \{D_{\ell} \mid \ell \in V\}$  be the set of all those double-rays (drawn thick in Figure 1). It has been shown in [13] that all the elements of  $\mathcal{D}$  and all the (continuum many) ends of F together form a circle C in |F| (the double rays  $D_l$  are traversed by C in the order they appear in Figure 1 when moving from left to right; one can think of C as being the boundary of the outer face of F).

The example we are going to use in order to see how the proof of Theorem 1.2 fails to prove Theorem 1.3 is the graph G obtained by taking two disjoint copies  $F_1, F_2$  of F and joining every vertex of  $F_1$  to its copy in  $F_2$  by an edge. The subgraph H of |G| we choose consists of the copies of  $\mathcal{D}$  in both  $F_1, F_2$ . Its closure  $\overline{H}$  contains all ends of G. Note that if R is a ray in  $F_1$  and R' its copy in  $F_2$ , then R and R' belong to the same end of G. This implies that G and Fhave the same ends and  $\Omega(G)$  and  $\Omega(F)$  have the same topology; formally, the function  $\pi : \Omega(F) \to \Omega(G)$  that maps every end of F to the end of G containing it as a subset is a homeomorphism. This means that the two copies  $C_1, C_2$  of C in  $F_1, F_2$  respectively are also circles in |G|, thus  $\overline{H}$ , which is the union of  $C_1$ and  $C_2$ , is path-connected since  $C_1$  and  $C_2$  intersect:  $C_1 \cap C_2 = \Omega(G)$ .

Thus H satisfies the conditions of Theorem 1.3: we have  $E(H) = E(C_1) \cup E(C_2) \in \mathcal{C}(G)$ , and  $\overline{H}$  is connected since it is path-connected. Now let us try to prove that H has a topological Euler tour in |G| using the proof of Theorem 1.2 sketched above. Firstly, we have to apply Theorem 1.5. If we are lucky this will decompose E(H) into  $E(C_1)$  and  $E(C_2)$ , and we will be able to imitate



Figure 1: The graph F. The thick double-rays together with the ends of the graph form a circle.

that proof (although we will have to attach the two circles at an end rather than at a vertex). However, we can be unlucky: Theorem 1.5 could return a quite different decomposition into circuits. Note that for every  $D \in \mathcal{D}$  the two copies of D in G together with their two incident ends form a circle  $C_D$  in |G|, comprising two double rays and two ends (recall that the two copies of a ray of F in G belong to the same end of G). Now E(H) can be decomposed into the set of all edge-sets of such circles  $C_D$  for all  $D \in \mathcal{D}$ . But if the application of Theorem 1.5 returns this decomposition then we cannot continue with the proof: if D, D' are any two distinct elements of  $\mathcal{D}$  then  $\overline{D} \cap \overline{D'} = \emptyset$ , so we cannot start with some circuit and then stepwise attach the others, since there are no common points where we could attach.

Thus we are going to need a different approach in order to prove our result:

Proof of Theorem 1.3. Let us start with the easier of the two implications: if H has a topological Euler tour in |G| then, since arcs cannot jump finite cuts by Lemma 3.3, every finite cut of G has an even number of edges in H. Lemma 3.2 now implies that  $E(H) \in C(G)$ .

We now assume conversely that  $E(H) \in \mathcal{C}(G)$  and construct a topological Euler tour  $\sigma$  of H in |G|.

Let  $v_1, v_2, \ldots$  be an enumeration of the vertices of V(G) and let  $G_n$  be the submultigraph  $G[v_1, v_2, \ldots, v_n]$  of G induced by the first n vertices in this enumeration. Moreover, let  $H_n := (V(G_n^*), E(H) \cap E(G_n^*))$  be the submultigraph of  $G_n^*$  formed by the edges in  $E(H) \cap E(G_n^*)$  and their incident vertices (the notation  $G_n^*$  is defined is Section 2).

Note that every cut of  $G_n^*$  is also a finite cut of G. Thus, if  $H_n$  intersects some cut of  $G_n^*$  in an odd number of edges, then so does H for the corresponding cut of G, which cannot be the case by Lemma 3.2. This means that for any bipartition  $\{X,Y\}$  of  $V(G_n^*)$ ,  $H_n$  contains an even number of X-Y edges. Moreover,  $H_n$  is connected, since otherwise there is a finite cut of  $G_n^*$ , and G, that separates  $H_n$  but contains no edge of H, and this cut easily contradicts the topological connectedness of  $\overline{H}$ . Thus by Euler's theorem [11]  $H_n$  has an Euler tour.

Note that every Euler tour W of  $H_{n+1}$  induces an Euler tour W|n of  $H_n$ ;

indeed, let W|n be the walk obtained from W by contracting all edges not in  $E(H_n)$ . By construction W|n contains all edges of  $H_n$  and traverses each of them only once.

We want to apply König's Infinity Lemma [23] (or [11]) to obtain a sequence  $(W_n)_{n\in\mathbb{N}}$  such that  $W_n$  is an Euler tour of  $H_n$  and  $W_{n+1}|_n = W_n$  for every n. So let  $V_n$  be the set of all Euler tours of  $H_n$ , and join any  $W \in V_{n+1}$  to  $W|_n \in V_n$ . The Infinity Lemma now yields a sequence  $(W_n)_{n\in\mathbb{N}}$  with  $W_n \in V_n$  as required ( $W_0$  being the trivial walk).

We are now going to construct a topological Euler tour  $\sigma$  of H in |G| as the limit of a sequence  $\sigma_n$  of topological Euler tours of the  $H_n$  corresponding to  $W_n$ .

Let  $\sigma_0$  be the constant map from  $S^1$  to the vertex  $u_C$  of  $G_0^*$  that corresponds to the component C of G that meets H. Assume inductively that for some nwe have defined  $\sigma_n : S^1 \to |G_n|$  so that  $\sigma_n$  runs through the edges and vertices of  $W_n$  in order, traversing every edge in a straight manner and pausing for a non-trivial interval at every vertex it visits (and each time it visits that vertex); more formally, the inverse image under  $\sigma_n$  of a vertex is always a disjoint union of non-trivial subarcs of  $S^1$ , one for each visit of  $W_n$  to that vertex.

We are now going to use  $\sigma_n$  in order to define  $\sigma_{n+1}$ . It is easy to see that  $W_n$  can be obtained from  $W_{n+1}$  by contracting every maximal subwalk that does not contain edges in  $W_n$  to a "contracted vertex"

$$u \in V(G_n^*) \backslash V(G_n).$$
(2)

For every such subwalk D, there is one subarc  $I_D$  of  $S^1$  mapped to u by  $\sigma_n$  that corresponds to D: there is an occurence of u in  $W_n$  that is split into two occurences of u in  $W_{n+1}$  between which D (and nothing else) is traversed, and  $I_D$  is the subarc that corresponds to that occurence.

Now modify  $\sigma_n$  so that it maps  $I_D$  continuously to D traversing every edge e of D precisely once, in a straight manner, and in the same direction as  $W_{n+1}$  traverses e, and pausing for a non-trivial interval at every vertex it visits. Let  $\sigma_{n+1}$  be the map obtained from  $\sigma_n$  by performing all these modifications for all contracted walks D (note that if D, D' are distinct maximal contracted walks then  $I_D \cap I_{D'} = \emptyset$ , thus all these modifications can be performed simultanously).

We now define the limit map  $\sigma: S^1 \to |G|$  as follows. Let  $x \in S^1$  be given. If there is a point y so that  $\sigma_n(x) = y$  for all but finitely many n, then put  $\sigma(x) = y$ . (In this case y must be a vertex of G or an inner point of an edge, and the maps  $\sigma_n$  will agree on x for all n so large that this vertex or edge is in  $G_n$ .) If not, then  $\sigma_n(x)$  is a "contracted vertex"  $u_{C_n}$  for every n by (2), and the corresponding components  $C_n$  are nested. Let  $\overline{C_n}$  denote the topological closure of  $C_n$  in |G|. As |G| is compact,  $U := \bigcap_n \overline{C_n}$  is non-empty, and as, clearly,  $\bigcap_n C_n = \emptyset$ , we have  $U \subseteq \Omega(G)$ . However, for any two distinct ends  $\omega, \omega'$  in  $\Omega(G)$  there is a finite vertex set S separating their rays. This means that if  $S \subseteq V(G_n)$  then  $\overline{C_n} = \omega$ . Note that

$$\sigma(x) \in \overline{C_n} \text{ for every } n. \tag{3}$$

This completes the definition of  $\sigma$  and now we have to show that it is continuous. This is trivial at points of  $S^1$  that are mapped to some vertex or inner point of an edge, so let x be a point mapped to an end  $\omega$ . We have to specify for every open neighbourhood O of  $\omega$  an open subset of  $S^1$  containing x mapped to O by  $\sigma$ . However, for every basic open set  $O' = O(S, \omega)$  of  $\omega$ , there is an  $n \in \mathbb{N}$  such that  $\overline{C} \subseteq O'$  for some component C of  $G - G_n$ ; just choose n large enough that  $S \subseteq \{v_1, \ldots, v_n\}$ . This component C corresponds to a vertex  $u_C \in V(G_n^*)$ , and the inverse image of  $u_C$  under  $\sigma_n$  is, by the construction of  $\sigma_n$ , a disjoint union I of non-trivial subarcs of  $S^1$ . We claim that  $\sigma(I) \subseteq \overline{C}$ . Indeed, if  $x \in I$  is mapped to a vertex or inner point of an edge by  $\sigma$ , then clearly  $\sigma(x)$  lies in C. If x is mapped to an end, then  $\sigma(x) \in \overline{C}$  holds by (3). This proves that  $\sigma$  is continuous. By construction  $\sigma$  traverses each edge in E(H) precisely once and traverses no other edge, thus it is a topological Euler tour of H in |G|.

### 5 An elementary proof of Theorem 1.5

In this section we prove Lemma 1.4 and apply it to obtain a proof of Theorem 1.5 which is shorter than previous proofs ([14, 11]).

Proof of Lemma 1.4. Let  $v_1, v_2, \ldots$  be an enumeration of the vertices of  $V(G) - \{x, y\}$  and let  $G_n$  be the submultigraph  $G[x, y, v_1, v_2, \ldots, v_n]$  of G induced by x, y and the first n vertices in this enumeration. Given any topological path  $\tau$  in |G| let  $\tau | n := \tau | G_n$ . We begin by selecting the sequence  $(\tau_{a_i})$ , in  $\omega$  steps, choosing one member in each step. We will choose  $(\tau_{a_i})$  so that all members with index i or greater will agree on  $G_i^*$ :

$$\tau_{a_i} | i = \tau_{a_i} | i \text{ for all } i, j \in \mathbb{N} \text{ with } i < j.$$

$$\tag{4}$$

For the first step, we pick an x-y trail  $W_1$  in  $G_1^*$  such that the set  $\mathcal{X}_1$  of indices j such that  $\tau_j|1 = W_1$  is infinite; such a trail exists because  $G_1^*$  is finite and  $\{\tau_i\}$  infinite. Then, pick any element j of  $\mathcal{X}_1$  and let  $\tau_{a_1} := \tau_j$ .

Having performed the first n-1 steps we pick, in step n, an x-y trail  $W_n$ in  $G_n^*$  such that the set  $\mathcal{X}_n$  of indices  $j \in \mathcal{X}_{n-1}$  such that  $\tau_j | n = W_n$  is infinite; such a trail exists because  $G_n^*$  is finite and  $\mathcal{X}_{n-1}$  infinite. Then, pick any element j of  $\mathcal{X}_n$  and let  $\tau_{a_n} := \tau_j$ .

Thus we have defined the sequence  $(\tau_{a_i})_{i \in \mathbb{N}}$ , which by construction satisfies (4). We will now construct a topological path  $\sigma$  from x to y in |G| with  $E(\sigma) \subseteq$  $\liminf(E(\tau_{a_i}))$  in  $\omega$  steps. For the first step, let  $\sigma_0$  be a continuous map from the real interval [0,1] to  $G_1^*$  that runs through each edge of  $W_1$  precisely once and in order, runs through no other edge, and maps a non-trivial interval  $I_v^j$  to each vertex v of  $G_1^*$  obtained by contracting a component of  $G - G_1$  each time it visits v.

Then, at step n, modify  $\sigma_{n-1}$  on these intervals  $I_v^j$  to extend it to a continuous map  $\sigma_n$  from [0, 1] to  $G_n^*$  that runs through each edge of  $W_n$  precisely once and in order, also mapping non-trivial intervals  $I_v^j$  to each contracted vertex vit visits; it is possible to extend  $\sigma_{n-1}$  in this way by (4).

We can now define the required topological path  $\sigma : [0,1] \to |G|$  as a "limit" of the  $\sigma_n$ : for every  $x \in [0,1]$ , if there is an  $n_0 \in \mathbb{N}$  such that  $\sigma_m(x) = \sigma_{n_0}(x)$ for every  $m \ge n_0$ , then we let  $\sigma(x) = \sigma_{n_0}(x)$ ; otherwise,  $\sigma_n(x)$  is a contracted vertex for every n, and as in the proof of Theorem 1.3 it is easy to prove that the components contracted to the  $\sigma_n(x)$  converge to a unique end  $\omega$  of G, and we let  $\sigma(x) = \omega$ . Similarly with the proof of Theorem 1.3, it is straightforward to check that  $\sigma$  is continuous, thus it is a topological x-y path. By construction  $\sigma$  has the required properties.

Note that Lemma 1.4 remains true if one or both of x, y is an end. To see this, modify the proof of Lemma 1.4 so that  $W_n$  is an  $\hat{x}-\hat{y}$  trail in  $G_n^*$ , where  $\hat{x}$  is the contracted vertex of  $G_n^*$  corresponding to the component containing rays in x if x is an end, and  $\hat{x} = x$  if x is a vertex; define  $\hat{y}$  similarly.

Proof of Theorem 1.5. It suffices to prove that every non-empty element of  $\mathcal{C}(G)$  contains at least one circuit; indeed, assume this is true, and for any  $C \in \mathcal{C}(G)$  let Z be a maximal set of edge-disjoint circuits in C. Then,  $C' = C + \sum_{D \in Z} D = C \setminus \bigcup_{D \in Z} D$  is an element of  $\mathcal{C}(G)$ , and if it is non-empty then it contains a circuit by our assumption, which contradicts the maximality of Z. Thus  $C' = \emptyset$  and C is the sum of the family Z.

To show that any non-empty  $C \in C(G)$  contains a circuit, we will pick an arbitrary edge  $f = xy \in C$  and find a circuit in C containing f. By definition, C is the sum of a thin family  $\{E(C_i)\}_{i\in\mathbb{N}}$  of circuits of circles  $C_i$ . Let  $F := \bigcup_{i\in\mathbb{N}} E(C_i)$ , let  $e_1, e_2, \ldots$  be an enumeration of the edges of G, and let  $E_n := \{f\} \cup \{e_i \mid i < n, e_i \in F \setminus C\}$ .

**Lemma 5.1.** For every  $n \in \mathbb{N}$  there is an x-y arc  $X_n$  in |G| with  $E(X) \subseteq F \setminus E_n$ .

*Proof.* Since  $\{C_i\}$  is a thin family, the subset  $K_n$  of its elements that contain edges in  $E_n$  is finite. For every  $C_i \in K_n$  replace every maximal subarc R of  $C_i$  that does not contain any edge in  $E_n$  by an edge  $e_R$ , called a *representing edge*, to obtain a finite cycle  $C'_i$ . Taking the union of all these new cycles we obtain a finite auxiliary graph

$$G_n := \left(\bigcup_{C_i \in K_n} V(C'_i), \bigcup_{C_i \in K_n} E(C'_i)\right)$$

and we can apply the finite version of Theorem 1.5 (see [11]) to the element  $\sum_{C_i \in K_n} C'_i$  of  $\mathcal{C}(G_n)$ . This yields a disjoint union of finite circuits, and one of them, say D, has to contain f since f must appear in an odd number of elements of  $\{C_i\}$  and thus of  $K_n$ . Moreover, D does not contain any edge in  $E_n - f$ , as those edges appear in an even number of elements of  $\{C_i\}$  and thus of  $K_n$ . Now replacing every representing edge  $e_R$  in D by the arc R it represents and removing f yields a topological x-y path which, by Lemma 3.1, contains the required x-y arc.

Applying Lemma 1.4 on the sequence  $(\tau_n)_{n \in \mathbb{N}}$  where  $\tau_n$  is a topological path whose image is the arc  $X_n$  obtained from Lemma 5.1 yields a topological x-y path all of whose edges lie in C, and applying Lemma 3.1 we obtain an x-y arc X. The union of X with f is the required circuit.

# 6 Cyclability

In this section we generalise Theorem 1.6 to the topological cycle space of a locally finite graph:

**Theorem 6.1.** If G is a locally finite multigraph and  $H \subseteq E(G)$  then H extends to an element of  $\mathcal{C}(G)$  if and only if H contains no cut of G of odd cardinality.

In order to prove the backward implication of Theorem 6.1 we will first show that it is true for finite H, and then use this fact and compactness to prove it for any H.

For H, G as in Theorem 6.1 let G[H] denote the submultigraph of G formed by the edges in H and their incident vertices.

**Lemma 6.2.** For every locally finite multigraph G and every finite  $H \subseteq E(G)$ , if H contains no cut of G of odd cardinality then it extends to an element of C(G).

*Proof.* We may assume without loss of generality that G is connected, since otherwise we can apply the result to each component of G that meets H.

We are going to find a finite submultigraph G' of G containing H such that H extends to an element of  $\mathcal{C}(G')$ . For this, consider every (possibly empty)  $F \subseteq E(H)$ , and if F separates G[H] but not G (in particular, if  $F = \emptyset$  and G[H] is not connected), then pick a set  $\mathcal{P}_F$  of paths in G so that  $(G[H] - F) \cup \bigcup \mathcal{P}_F$  is connected. Let G' be the (finite) union of G[H] with all paths that lie in some  $\mathcal{P}_F$ . By the construction of G', if H contains a cut B of G' then B is also a cut of G. Thus H contains no cut of G' of odd cardinality, and we can apply Theorem 1.6 to H, G' to prove that H extends to an element of  $\mathcal{C}(G')$ . Since any cycle in G' is also a cycle in G, this means that H also extends to an element of  $\mathcal{C}(G)$ .

We can now prove Theorem 6.1 applying Lemma 6.2 to the finite subsets of H and using compactness. One way to do this is using the methods of the previous sections and König's Infinity Lemma as the reader can check, however the most convenient way is to use the compactness theorem for propositional logic (see [6]):

**Theorem 6.3.** Let K be an infinite set of propositional formulas every finite subset of which is satisfiable. Then K is satisfiable.

Proof of Theorem 6.1. If H contains a cut of G of odd cardinality then by Lemma 3.2 it does not extend to an element of  $\mathcal{C}(G)$ .

For the reverse implication, let  $e_1, e_2, \ldots$  be an enumeration of the elements of H, and for every  $i \in \mathbb{N}$  let  $H_i = H \cap \{e_1, \ldots, e_i\}$ . By Lemma 6.2 there is for each  $i \in C(G)$  with  $H_i \subseteq C_i$ . In order to apply Theorem 6.3, we introduce for each edge e of G a boolean variable  $v_e$ ; intuitively, the two possible values T, F that  $v_e$  can assume encode the presence or not of e. For every edge e in H let  $P_e$  denote the propositional formula  $v_e = T$  (encoding the event that e is present). Moreover, for every finite cut B of G let  $P_B$  be a propositional formula (in the variables  $v_e$ ) that evaluates to T if and only if the set  $\{e \in B \mid v_e = T\}$ has even cardinality ( $P_B$  expresses the event that an even number of edges in B are present). Let  $K := \{P_e \mid e \in H\} \cup \{P_B \mid B \text{ is a finite cut of } G\}$  be the set of all these propositional formulas. We claim that every finite subset K' of K is satisfiable. Indeed, pick  $n \in \mathbb{N}$  large enough that  $H_n$  contains all edges e such that  $P_e \in K'$ . Now  $C_n$  yields an assignment of values to the variables  $v_e$  that satisfies all formulas in K': for every  $e \in E(G)$  let  $v_e = T$  if  $e \in C_n$  and  $v_e = F$  otherwise. It is obvious that this assignment satisfies all formulas in K' of the form  $P_e$ . It also satisfies all formulas of the form  $P_B$  as, by Lemma 3.2,  $C_i$  meets every finite cut in an even number of edges.

We can thus apply Theorem 6.3 to obtain an assignment of values to the  $v_e$  that satisfies all formulas in K. The corresponding edge set  $C := \{e \in E(G) \mid v_e = T\}$  contains, then, all edges in H, and it lies in  $\mathcal{C}(G)$  by Lemma 3.2.

Theorem 1.6 is proved by an elegant algebraic argument ([22]), and in fact this argument can be adapted to give an alternative proof to Theorem 6.1. Let G be a finite multigraph. For a subset F of E(G), denote by P(F) the subspace of the edge space  $\mathcal{E}$  of G generated by the edges in F, and let  $\overline{F} := E(G) \setminus F$ . For  $X \subseteq \mathcal{E}$  denote by  $X^{\perp}$  the set of elements of  $\mathcal{E}$  that meet each element of Xin an even number of edges. Moreover, let K be the subspace of  $\mathcal{E}$  consisting of the cuts of G, and let  $\mathcal{C}$  be the cycle space of G. In this notation, Theorem 1.6 can be formulated as follows:

If G is a finite multigraph and 
$$H \subseteq E(G)$$
 then  $H \in \mathcal{C} + P(\overline{H})$   
if and only if  $H \in [K \cap P(H)]^{\perp}$ . (5)

Now to prove (5) note that  $[\mathcal{C} + P(\bar{H})]^{\perp} = \mathcal{C}^{\perp} \cap P(\bar{H})^{\perp} = K \cap P(H)^{\perp}$ , thus (5) follows from the fact that  $X^{\perp \perp} = X$  (it is well-known that  $\mathcal{C}^{\perp} = K$ ; see [11]). In order to apply the same argument in the case that G is an infinite locally finite multigraph, we need three facts (if X is a subspace of the edge space  $\mathcal{E}$  of G then  $X^{\perp}$  consists of those elements of  $\mathcal{E}$  that have a (finite and) even intersection with each element of X). Firstly, we may assume without loss of generality that G is 2-edge-connected, and thus Lemma 3.2 implies that  $\mathcal{C}^{\perp}$ is the space of finite cuts of G. Second, that  $\mathcal{C}(G)$  is generated by a thin set (see again [11, Theorem 8.5.8]), and finally, that if  $X \subseteq \mathcal{E}$  is generated by a thin set then  $X^{\perp \perp} = X$ , which follows from Theorem 4.3 of [7].

Let me now argue that Theorem 6.1 could find applications in the study of hamiltonicity in infinite graphs. If G is a locally finite graph then we define a *Hamilton circle* of G to be a circle in |G| containing all vertices (and hence also all ends, as it is closed). This concept allowed the generalisation of Fleischner's theorem to locally finite graphs [18]. See [19] for more results and open problems about Hamilton circles.

Thomassen [24] conjectured that every finite 4-connected line graph is hamiltonian. An easy special case of this conjecture is to prove that the line graph of a 4-edge-connected finite graph G is hamiltonian. To prove this, note that since G is 4-edge-connected it has two edge-disjoint spanning trees T, S ([11, Theorem 2.4.1]). By Theorem 1.6, E(T) extends to an element C of C(G): it contains no cut of G since its complement contains the spanning tree S. So let W be an Euler tour of G[C], which exists by Theorem 1.2. Since W meets all vertices of G it is easy to convert W into a Hamilton cycle of the line graph L(G) of G: we just translate edges of W into vertices of L(G), and each time we visit a new vertex v of G while running along W we make a detour to pick up any edges incident with v that are not in W and have not been picked up yet; see [8] for a more precise exposition. Having generalised Theorem 1.6 to Theorem 6.1 we could try to extend this proof to locally-finite graphs G. However the first step is still missing: we need a spanning subgraph of G that has the same end-topology as G and whose edge-set extends to an element of  $\mathcal{C}(G)$  to play the role of T in the above sketch.

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